Subregular toolkit implemented in Python

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Subregular toolkit: general information

kist: kist implementing subregular toolkit

Motivation: to collect in one place the functionality for subregular languages and subsequential transducers.

- For researchers: to avoid manual burden of extracting grammars and designing transducers, creating data samples, or scanning strings;
- For practitioners: to start using tools in practice that are currently available only in the literature.
- Python 3 (will be available via pip)
- Open source
- Available on GitHub
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More motivations

This subregular toolkit allows one to:

- use recent theoretical results in practice;
- test ideas currently available in the literature;
- explore new methods to model natural language;
- automatically extract dependencies therefore avoiding manual burden of automata/transducer construction.



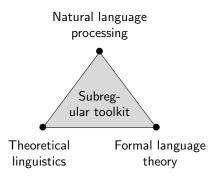
More motivations

Introduction

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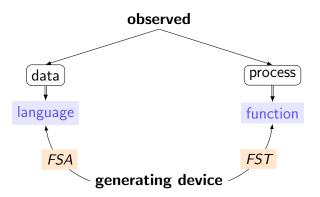
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The importance of formalization

- In order to abstract away from details and look at the big picture, we need to *formalize*:
 - Languages → sets of strings of a particular type;
 - Functions → descriptions of processes.
- KIST toolkit provides functionality that allows one to work with (sub)regular languages and functions.
- Such a toolkit is useful for NLP, and not only.

Languages vs. Functions



Here, I only work with (sub)regular – requiring a finite amount of memory – languages and functions.

What is done and what is left

Introduction

Last year:

- ✓ FSA implementation:
 - ✓ architecture;
 - optimization.
- ✓ Languages (SL, TSL, SP):
 - ✓ learners;

 - corresponding FSA.

What is done and what is left

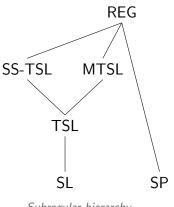
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 - ✓ scanners:
 - ✓ sample generators;
 - ✓ neg⇔pos switch;
 - corresponding FSA.

This year:

- □ Languages (MTSL, SS-TSL):
 - □ learners;
 - □ scanners;
 - □ sample generators;
 - □ neg↔pos switch;
 - □ corresponding FSA.
- □ Transduction learners:
 - □ OSTIA;
 - □ ISLFLA;
 - □ OSLFIA.

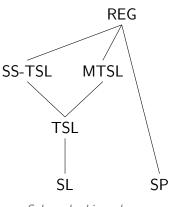
Languages and FSMs



Subregular hierarchy (simplified)

- The class of regular languages consists of smaller sub-classes. (McNaughton&Papert 1971)
- For every (sub)regular language, it is possible to construct a corresponding finite state automaton.
- Most subregular classes are learnable in polynomial time with positive data only.
- There is a variety of applications for subregular languages!

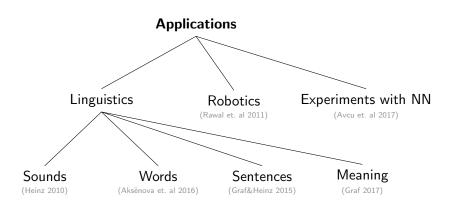
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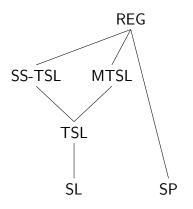
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What are the applications?



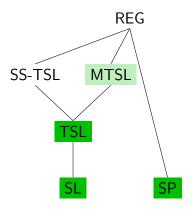
Subregular languages in KIST



Implemented functionality:

- learners;
- scanners;
- sample generators;
- negative ↔ positive grammar translators;
- constructing corresponding FSA;
- trimming FSA.

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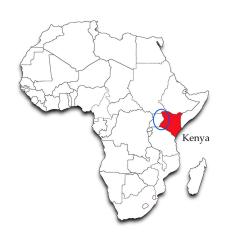
Language example

- Language: BUKUSU (Kenya)
- Construction: V + el/er/il/ir 'use something to V'
- Rule: "match the sounds of the suffix with the sounds of the verb"
- tleex-el 'use smth to cook'
- reeb-er 'use smth to ask'
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Simple formal version of the pattern: $(1,e)^+ \cup (1,i)^+ \cup (r,e)^+ \cup (r,i)^+$

ok Illiiilliiii

liiirriii

• ok eeerreer

• Teeelliii

• ok riiriirrr

• ...

Intuition is that [e] and [i] need to agree with each other, as well as [1] and [r]. Among themselves, these two agreements do not interact.

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ok |||iii|||iiii

• "liiirriii

• ok eeerreer

• oklleeelle

• ...

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$$(1,e)^+ \cup (1,i)^+ \cup (r,e)^+ \cup (r,i)^+$$

- Complexity: MTSL (multiple tier-based strictly local)
- *Meaning:* there are several sets of items involved in long-distance dependency.

•
$$T_1 = \{I, r\}$$
, and $G_{1_{pos}} = \langle II, rr \rangle$

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$$T_2 = \{e, i\}$$
, and $G_{2_{pos}} = \langle ee, ii \rangle$





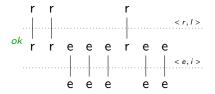
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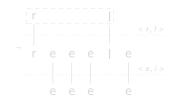
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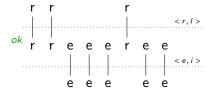
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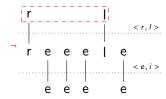




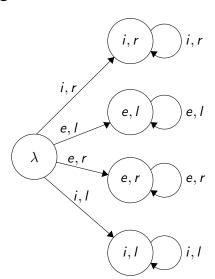
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Corresponding FSA:



Languages in kist: outcomes

- Aksënova, Alëna and Sanket Deshmukh (2018) Formal Restrictions on Multiple Tiers Proceedings of SCiL-2018, ACL anthology, Salt Lake City.
- Aksënova, Alëna (2018) The Hitchhiker's Guide to Harmony Interactions Poster at GLOW41, Budapest.
- McMullin, Kevin, Alëna Aksënova and Aniello De Santo (submitted) Learning Phonotactic Restrictions on Multiple Tiers
- Moradi, Sedigheh, Alëna Aksënova and Thomas Graf (submitted) The Computational Cost of Explicit Generalizations







Sanket Deshmukh



Kevin McMullin



Sedigheh Moradi

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- They are learnable from positive data only.
- For every subregular pattern, it is possible to construct a FSA.

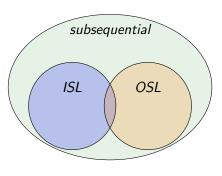
- A Finite State Automaton detects whether a given string belongs to a certain class.
- In order to re-write a string, one needs a **Finite State Transducer**.

Languages: local summary

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Functions and string FSTs



- String transducers have been used for different tasks since 1960-s. (Schützenberger 1961)
- In linguistics, multiple string extending and rewriting operations are represented via transductions. cat + s → cats

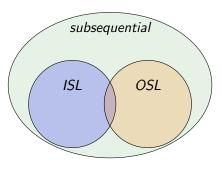
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(Chandlee 2014, i.a.)

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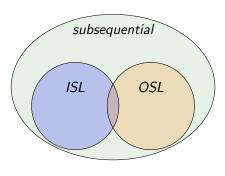
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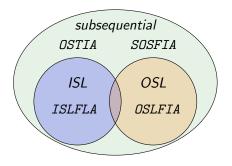
Learners for string transducers

- There are numerous learners for transductions. Among them:
 - OSTIA: subsequential transductions, cubic time, less data (Oncina, García, and Vidal 1993; de la Higuera 2010)
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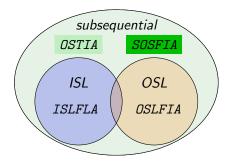
Subsequential transducers in KIST



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- transducer's template construction;
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String FSTs: an example of application

 Tokenization – separating words from sentence-level punctuations for further sentence processing.

```
"Bob, Sue and Bill didn't buy sugar-free coffee."
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- $\,\mapsto\,$ " Bob , Sue and Bill didn't buy sugar-free coffee ."
- Challenges:
 - Trying to avoid hard-coding the linguistic variety of contexts and punctuations (for example, Spanish '¿')
 - Not all same symbols are treated in the same way:
 "Dogs bark." → " Dogs bark ."
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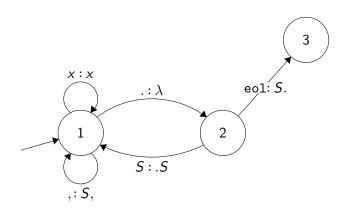
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String FSTs: an example of application [cont.]

Simplified FST for tokenization:



Functions in kist: current projects

 with Jeffrey Heinz and Kyle Gorman OSTIA-based tokenizer developing a low memory resource and high-accuracy tokenizer that avoids hard-coding language-specific information



Kyle Gorman

with Thomas Graf and Jeffrey Heinz
 Transduction learner for insufficient data
 creating a learning algorithm that allows to learn a class of non-equivalent transducers that can be inferred based on the insufficient input data

String transducers: local summary

- Finite State Transducers read a string as input, and return another string as output.
- Variety of different tasks can be performed via FSTs:
 - tokenization;
 - XML parsing;
 - multiple linguistic processes;
 - even machine translation!
- Current lines of research:
 - tree transducers;
 - one-to-many transductions;
 - learning of equivalent transducers for insufficient data;
 - ...and many others.

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Timeline

| Months | Goals |
|-----------|--|
| September | OSTIA (general subsequential learner) |
| October | OSLFIA (OSL transductions learner) |
| November | ISLFLA (ISL transductions learner) |
| December | MTSL learners, scanners, sample generators |
| January | SS-TSL learners, scanners, sample generators |
| February | Learner for regular languages (RPNI) |
| March | |
| April | Testing, documentation and publishing |
| May | |

Conclusion

- kist package is a subregular toolkit for linguistics and NLP.
- Last year, I implemented subregular language tools.
 Why? They learn and generate formal languages of a required complexity.
- This year, I am implementing transduction learners.
 Why? They extract different types of maps from input to output forms.
- For researchers, this toolkit facilitates the process of data analysis and generation, as well as assists in measuring complexities of already existent datasets.
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What I cannot create, I do not understand.

Richard Feynman

Thank you!



References I



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